

On the Complexity of Sampling Redistricting Plans

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Based on joint work with Moses Charikar, Paul Liu and Tianyu Liu

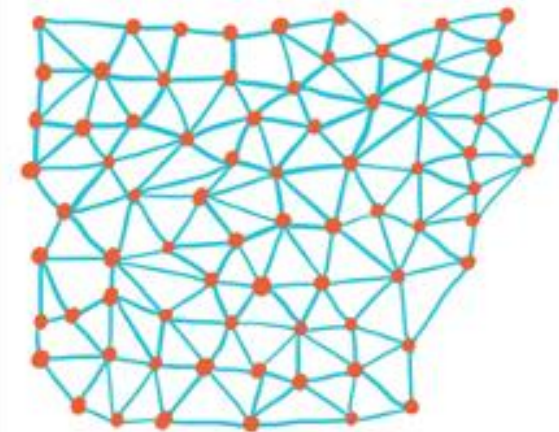
Outline

1. Redistricting problem
2. Sampling redistricting plans using Markov chains
 - Spanning tree distribution
 - Recombination (ReCom) chain
3. Slow mixing of ReCom on natural subgraph of the grid
4. Fixing ReCom
 - Fast mixing of relaxed ReCom
 - Enforcing balanced constraint

Redistricting

Partitioning map into districts = partitioning vertices of dual graph

Property of graph: planar, low degree (think: grid/subgraph of grid)



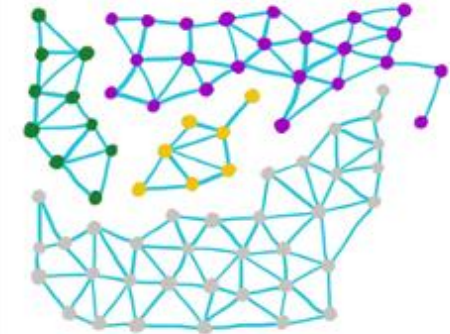
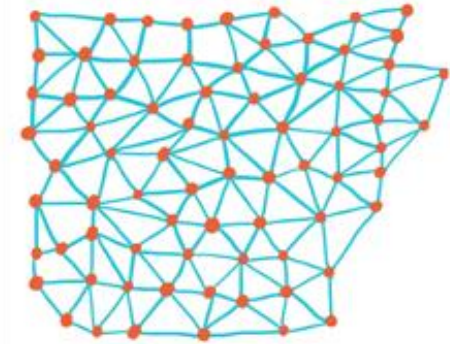
Redistricting

Partitioning map into districts = partitioning vertices of dual graph

Desirable properties for partitioning

$$P_1, \dots, P_k \subseteq V$$

- Balanced: equal population $|P_i| \approx |V|/k$
- Contiguous: $G[P_i]$ is connected
- Compact: few edges between different parts (small cut)



Gerrymandering

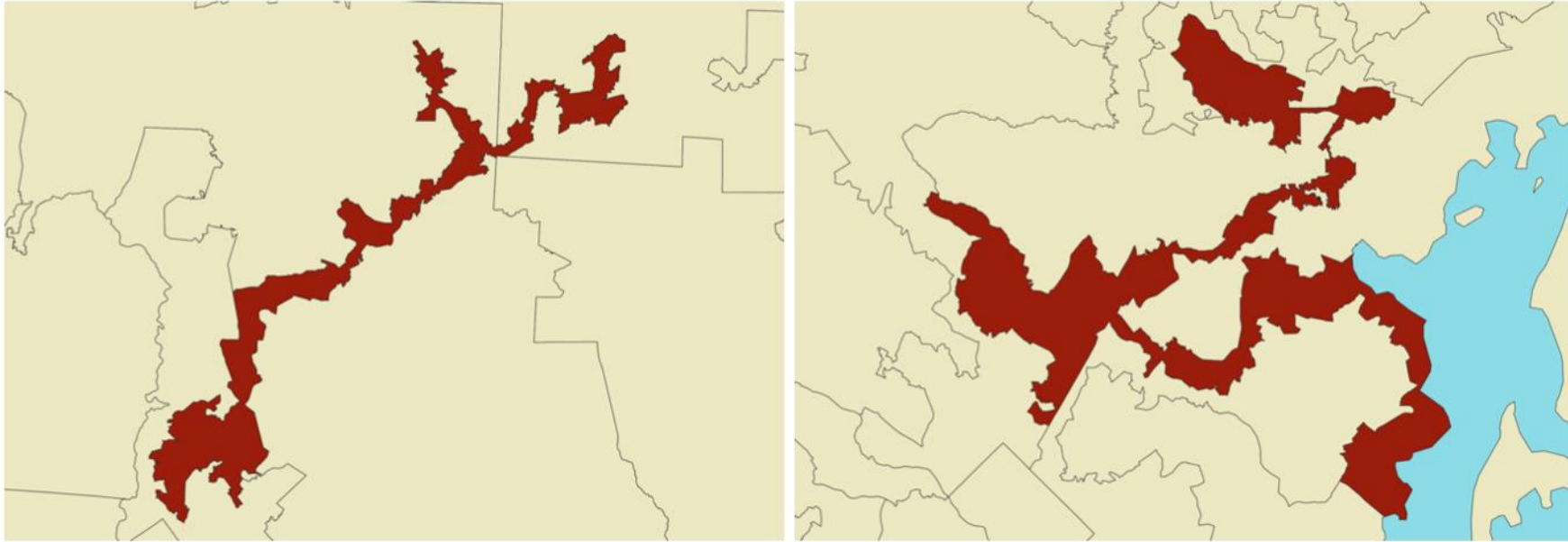


Figure 4: Examples of scapegoating: North Carolina's 12th district and Maryland's 3rd, circa 2013.

Detecting gerrymandering

Is a given redistricting plan $P = (P_1, \dots, P_k)$ an instance of gerrymandering?



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Solution: **sample** many redistricting plans from some bipartisan agreed-upon distribution, then show that P is “similar” to these plans

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Which distribution to sample from?

Goal: distribution that favor balanced, contiguous, compact partitioning

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But sampling from this is NP-hard ☹️

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Idea 1: partition with few cut edges: $\mu(P) \propto \lambda^{\#cut-edges(P)}$

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Idea 2: few cut edges = many spanning tree in each part

[Procaccia—Tucker-Foltz—SODA'21]

Spanning tree distribution

Let $T(P_i) = \# \{\text{spanning trees in } G[P_i]\}$.

Distribution μ over k-partition of graph G s.t.

$$\mu(P_1, \dots, P_k) \propto \prod T(P_i)$$

$\mu^{balanced}$: μ restricted to balanced partitioning

ReCom chain [Deford-Duchin-Solomon'19]

Current partition (P_1, \dots, P_k)

- Merge step: merge two parts P_i, P_j
- Split step: split the union to P'_i, P'_j

Can choose the probability of merge and split so that the Markov chain converges to $\pi \equiv \mu^{balanced}$

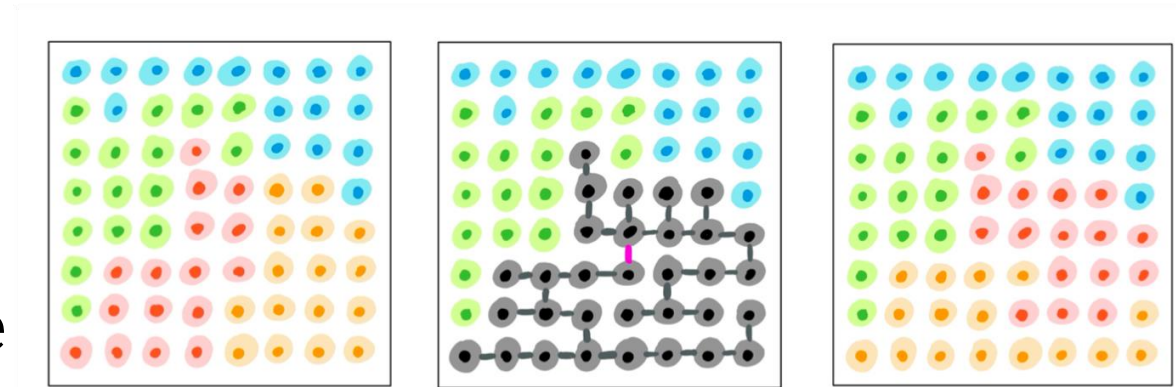
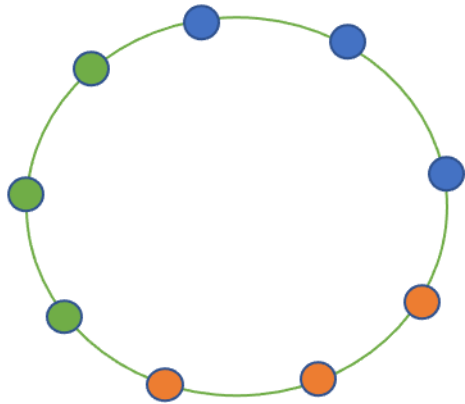


Figure 11: The basic recombination step: two districts are merged, a spanning tree is chosen, a balance edge is selected and cut, leaving two new districts.

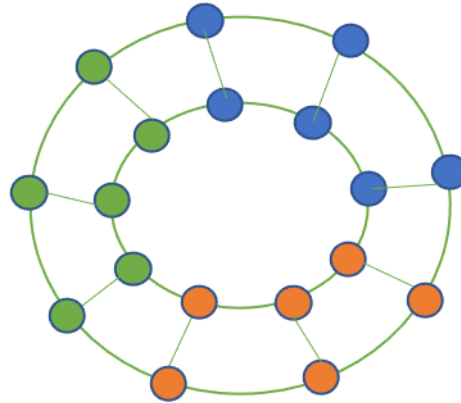
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 - c -weighted spanning tree distribution
 - Proportion of balanced partitionings

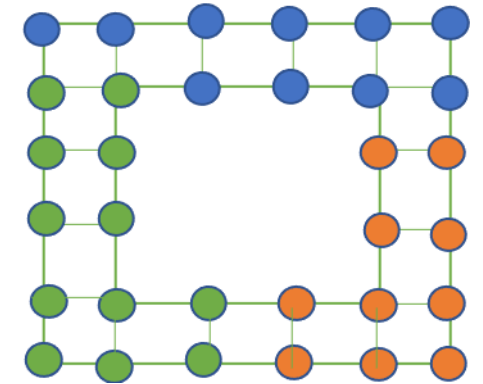
Slow mixing of ReCom



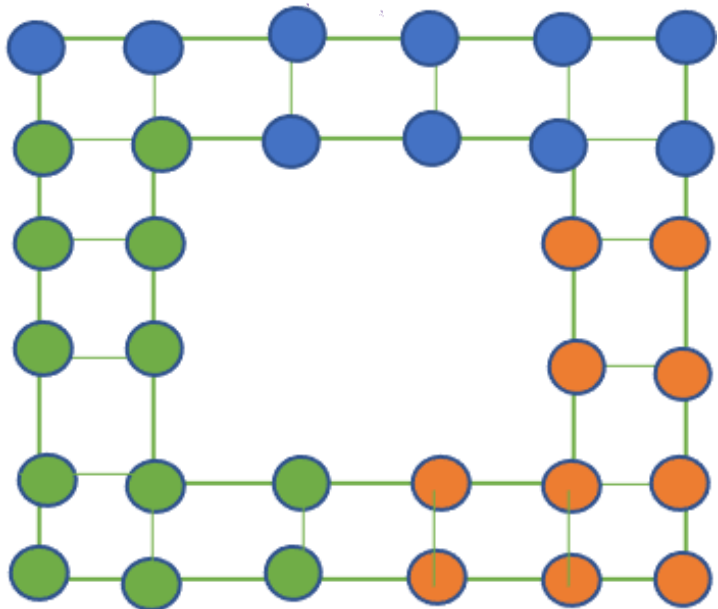
ReCom is not connected
on the single cycle



ReCom is connected on
the double cycle, but
takes $\exp(\Omega(n))$ to mix



ReCom is connected on
grid-with-a-hole, but
takes $\exp(\Omega(n))$ to mix



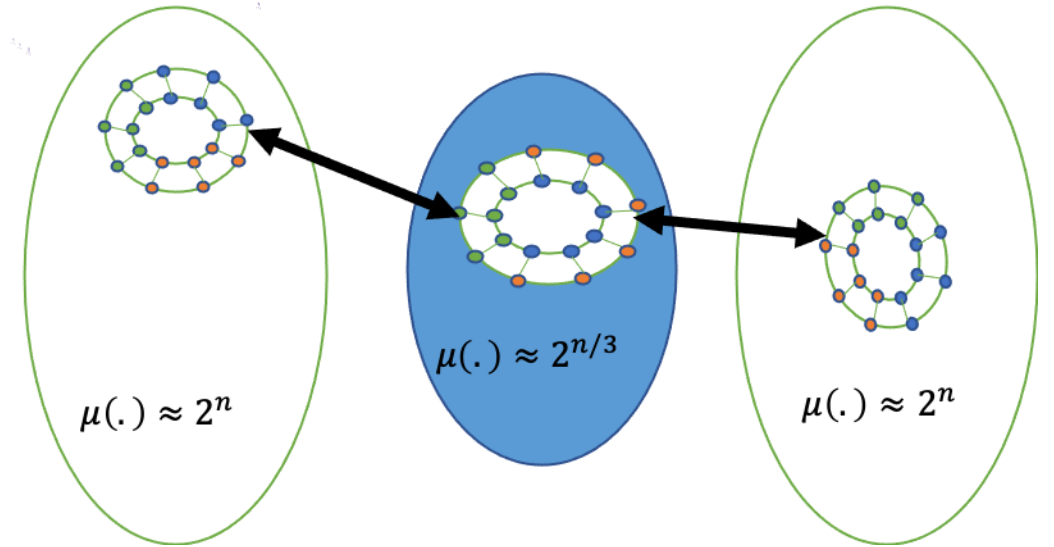
Natural example of grid-with-a-hole: map around a lake

Double cycle: proof intuition

Show existence of small cutset:

A set C of states (partitionings)
whose removal disconnect two states
 A and B s.t.

$$\frac{\pi(A)}{\pi(C)} \approx \frac{\pi(B)}{\pi(c)} = \exp(\Omega(n))$$



Double cycle: proof intuition

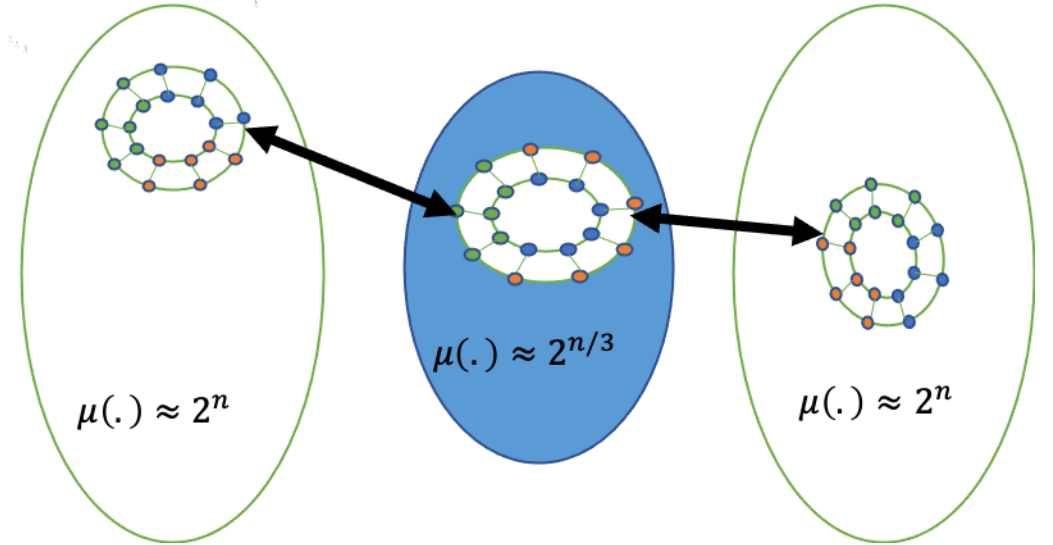
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Implies conductance $< \exp(-\Omega(n))$

thus slow mixing



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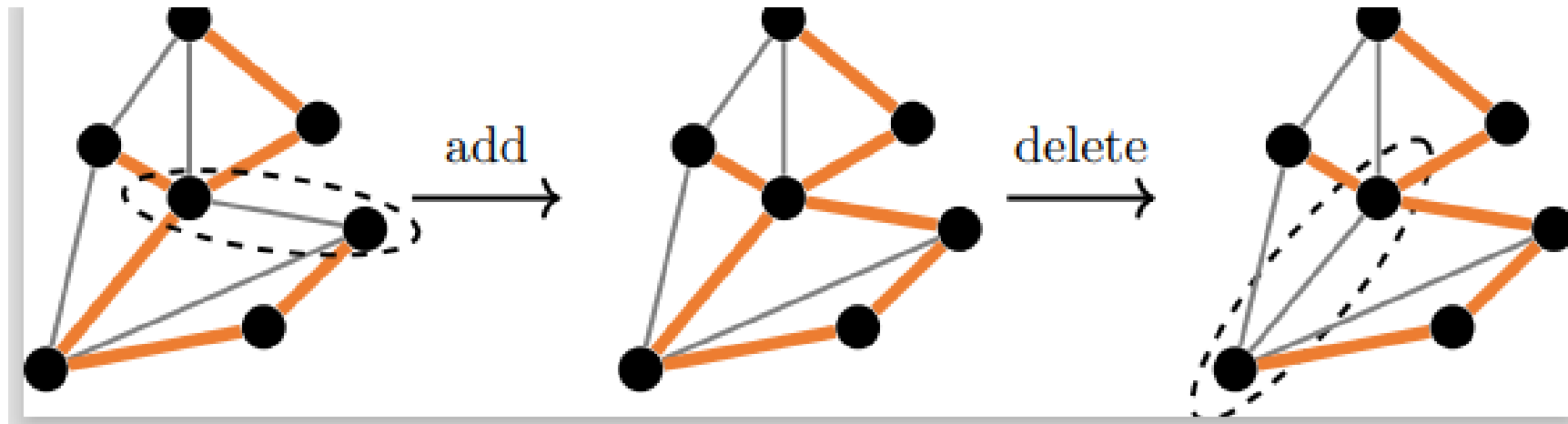
“Relaxed”-ReCom on forests

Current state (F_1, \dots, F_k) where F_i is a tree in $G[P_i]$, $P_i = V(F_i)$

- Merge step: merge two trees F_i, F_j by adding an edge
- Split step: split into two new trees F'_i, F'_j by removing an edge

This is precisely “Up-down walk on forest” [Anari-Liu-OveisGharan-Vinzant-Vuong-STOC'21] and runs in $O(|E| \log |E|)$ time

“Relaxed”-ReCom on forests



Enforcing balanced constraint

Q: How to enforce balanced constraint?

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A: Rejection sampling: sample from μ and accept only balanced partitionings

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On $m \times n$ grid-graph, experimental evidence that for $k = O(1)$

$$\Pr_{\mu}[\text{balanced}] = 1/\text{poly}(m, n)$$

Proof for the case $m = O(1)$, for double cycle/grid-with-a-hole.

Soft balanced constraint

C-weighted spanning tree distribution: μ^c over k-partition of graph G
s.t.

$$\mu^c(P_1, \dots, P_k) \propto \prod T(P_i) |P_i|^c$$

$$\mu^c = \mu \text{ if } c=0 \quad \mu^c = \mu \text{ if } c=0$$

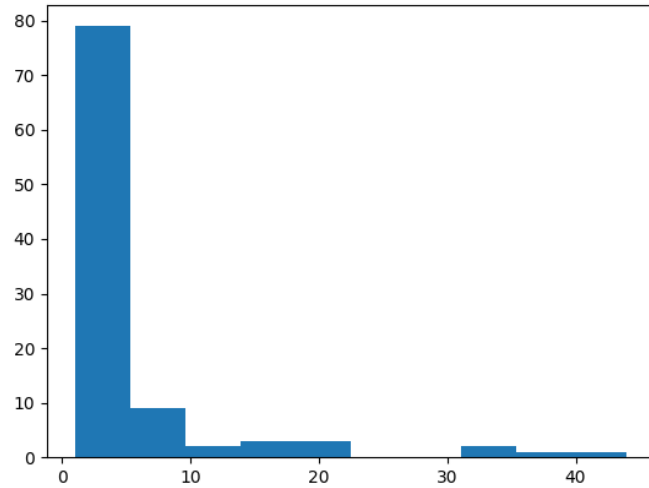
$$\mu^c = \mu^{balanced} \text{ if } c = \infty$$

Soft balanced constraint

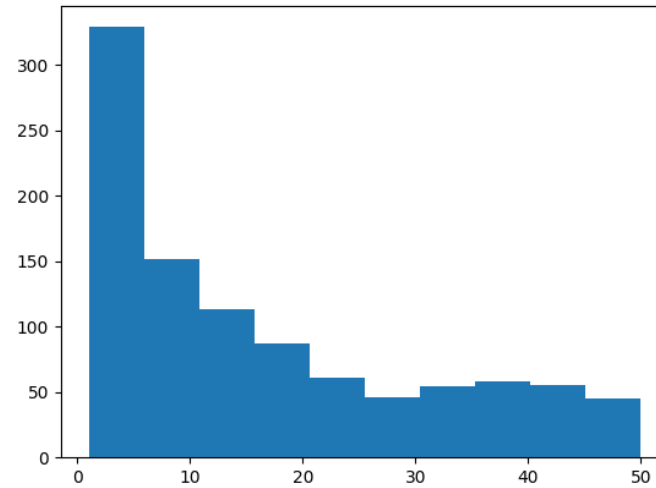
Similar up-down walk can sample from μ^c in time $\tilde{O}(|V|^{O(c)}|E|)$

(Analysis using comparison with μ^0)

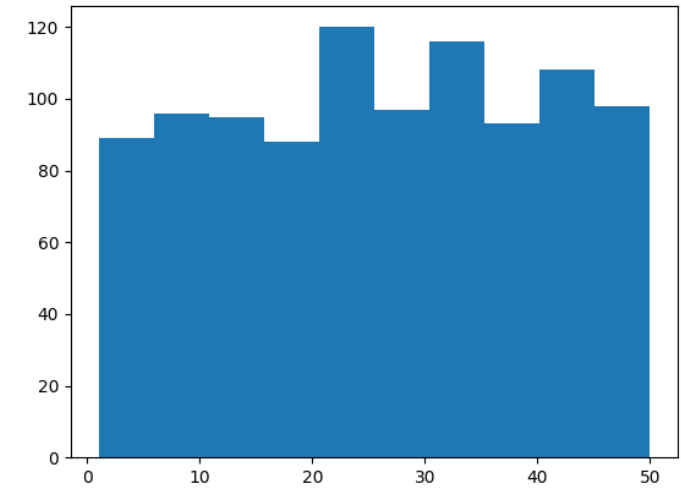
Proportion of balanced constraint



C=0



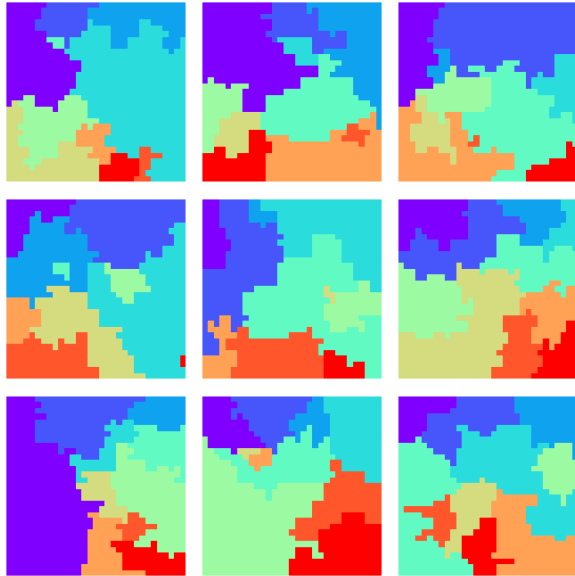
C=1



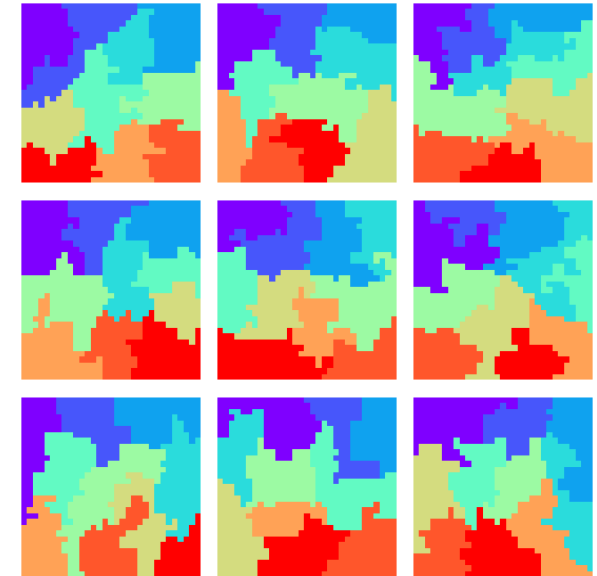
C=2

10x10 grid graph, $k=2$, histogram for size of minimum connected component

Proportion of balanced constraint



$C=2$



$C=20$

30x30 grid graph, $k=10$, example of sampled partitionings for different values of C